

Stony Brook University CSE 306: Operating Systems

Operating Systems History and Overview

Portions of this material courtesy Profs. Wong and Stark

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So what is an OS?

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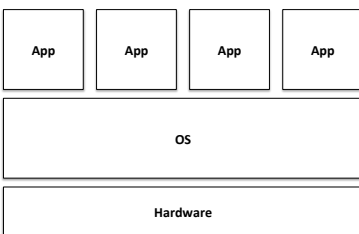
One view of an OS



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Another simple view of an OS



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A less happy view of an OS



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So which one is right?

- They all are

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An OS serves three masters

1. Give users a desktop environment
2. Give applications a more usable abstraction of the hardware
3. Give hardware manufacturers an abstraction of the applications

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Background (1)

- CPUs have 2 modes: user and supervisor
 - Sometimes more, but whatever
- Supervisor mode:
 - Issue commands to hardware devices
 - Power off, Reboot, Suspend
 - Launch missiles, Do awesome stuff
- User mode:
 - Run other code, hardware tattles if you try anything reserved for the supervisor

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OS architecture

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OS architecture

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Master #2: Applications

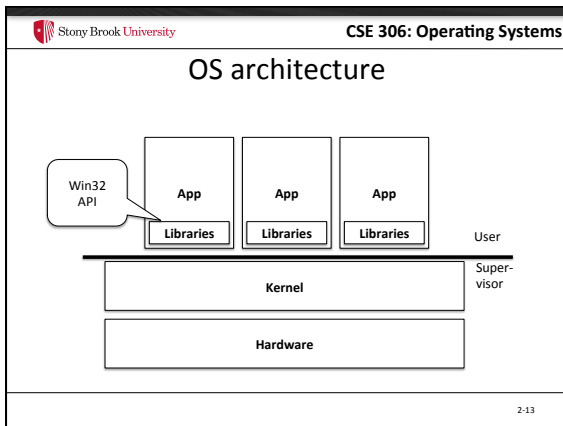
- Application Programming Interface (API)
 - Win32 (Windows)
 - POSIX (Unix/Linux)
 - Cocoa/Cocoa Touch (Mac OS/iOS)
- Application-facing functions provided by libraries
 - Injected by the OS into each application

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OS architecture

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Famous libraries, anyone?

- Windows: ntdll.dll, kernel32.dll, user32.dll, gdi32.dll
- Linux/Unix: libc.so, ld.so, libpthread.so, libm.so

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Caveat 1

- Libraries include a lot of code for common functions
 - Why bother reimplementing sqrt?
- They also give high-level abstractions of hardware
 - Files, printer, dancing Homer Simpson USB doll
- How does this work?

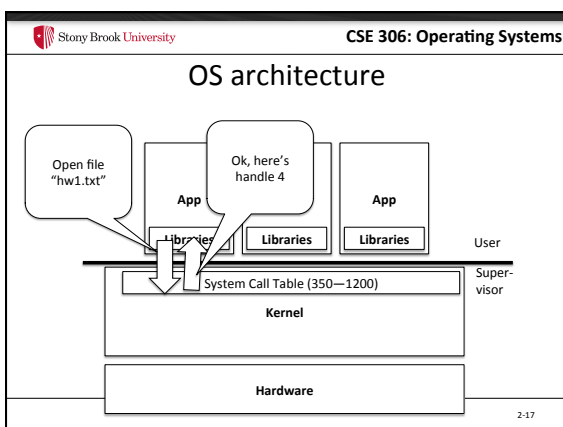
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System Call

- Special instruction to switch from user to supervisor mode
- Transfers CPU control to the kernel
 - One of a small-ish number of well-defined functions
- How many system calls does Windows or Linux have?
 - Windows ~1200
 - Linux ~350

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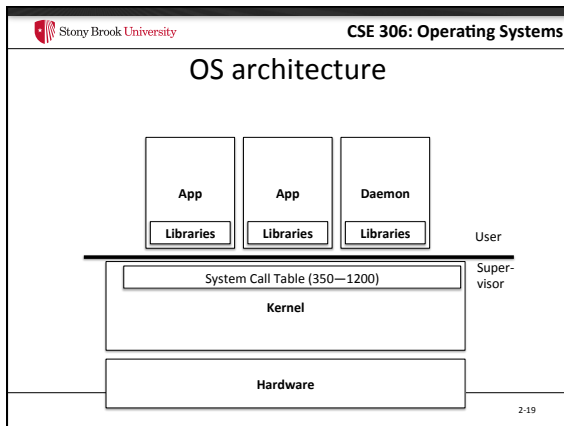


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Caveat 2

- Some libraries also call special apps provided by the OS, called a **daemon (or service)**
 - Communicate through kernel-provided API
- Example: Print spooler
 - App sends pdf to spooler
 - Spooler checks quotas, etc.
 - Turns pdf into printer-specific format
 - Sends reformatted document to device via OS kernel

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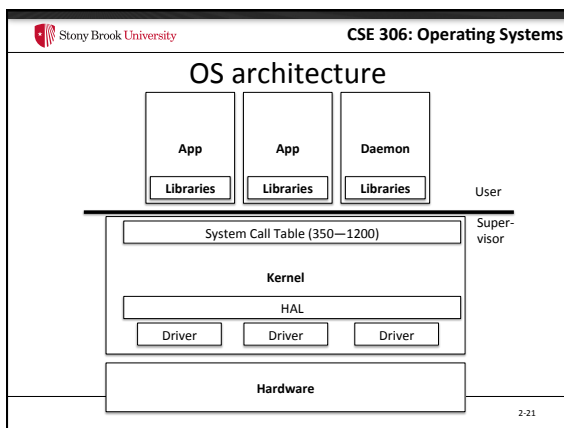


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Master 3: Hardware

- OS kernels are programmed at a higher low level of abstraction
 - Disk blocks vs. specific types of disks
- For most types of hardware, the kernel has a “lowest common denominator” interface
 - E.g., Disks, video cards, network cards, keyboard
 - Think Java abstract class
 - Sometimes called a hardware abstraction layer (HAL)
- Each specific device (Nvidia GeForce 600) needs to implement the abstract class
 - Each implementation is called a **device driver**

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What about Master 1

- What is the desktop?
- Really just a special daemon that interacts closely with keyboard, mouse, and display drivers
 - Launches programs when you double click, etc.
 - Some program libraries call desktop daemon to render content, etc.

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An OS serves three masters

- Give users a desktop environment
 - Desktop, or window manager, or GUI
- Give applications a more usable abstraction of the hardware
 - Libraries (+ system calls and daemons)
- Give hardware manufacturers an abstraction of the applications
 - Device Driver API (or HAL)

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Multiplexing Resources

- Many applications may need to share the hardware
- Different strategies based on the device:
 - Time sharing: CPUs, disk arm
 - Each app gets the resource for a while and passes it on
 - Space sharing: RAM, disk space
 - Each app gets part of the resource all the time
 - Exclusive use: mouse, keyboard, video card
 - One app has exclusive use for an indefinite period

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So what is Linux?

- Really just an OS kernel
 - Including lots of device drivers
- Conflated with environment consisting of:
 - Linux kernel
 - Gnu libc
 - X window manager daemon
 - CUPS printer manager
 - Etc.

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So what is Ubuntu? Centos?

- A **distribution**: bundles all of that stuff together
 - Pick versions that are tested to work together
 - Usually also includes a software update system

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OSX vs iOS?

- Same basic kernel (a few different compile options)
- Different window manager and libraries

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What is Unix?

- A very old OS (1970s), innovative, still in use
- Innovations:
 - Kernel written in C (first one not in assembly)
 - Co-designed C language with Unix
 - Several nice API abstractions
 - Fork, pipes, everything a file
- Several implementations: *BSDs, Solaris, etc.
 - Linux is a Unix-like kernel

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What is POSIX?


- A standard for Unix compatibility
- Even Windows is POSIX compliant!

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History of Operating Systems

- Two ways to look at history:
 - Evolution of the Theory
 - Evolution of the Machine/Hardware



The image shows three logos side-by-side: the Apple logo, the Windows logo, and the Ubuntu logo.

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Evolution of OS Theory

1. Centralized operating system
 - Resource management and multiprogramming, *Virtuality*
2. Network operating system
 - Resource sharing to achieve *Interoperability*
3. Distributed operating system
 - Single computer view of a multiple computer system for *Transparency*
4. Cooperative autonomous system
 - Cooperative work with *Autonomicity*

Decreasing Degree of Hardware and Software Coupling

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Evolution of OS Machine/Hardware

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1940's – First Computers

- One user/programmer at a time (serial)
 - Program loaded manually using switches
 - Debug using the console lights
- ENIAC
 - 1st gen purpose machine
 - Calculations for Army
 - Each panel had specific function

ENIAC (Electronic Number Integrator and Computer)

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1940's – First Computers

- Vacuum Tubes and Plugboards
- Single group of people designed, built, programmed, operated and maintained each machine
- No Programming language, only absolute machine language (101010)
- O/S? What is an O/S?
- All programs basically did numerical calculations

Pros:

- Interactive – immediate response on lights
- Programmers were women ☺

Cons:

- Lots of Idle time
 - Expensive computation
- Error-prone/tedious
- Each program needs all driver code

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1950's – Batch Processing

- Deck of cards to describe *job*
- Jobs submitted by multiple users are sequenced automatically by a *resident monitor*
- *Resident monitor* was a basic O/S
 - S/W controls sequence of events
 - Command processor
 - Protection from bugs (eventually)
 - Device drivers

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Monitor's Perspective

- Monitor controls the sequence of events
- *Resident Monitor* is software always in memory
- Monitor reads in job and gives control
- Job returns control to monitor

Figure 2.3 Memory Layout for a Resident Monitor

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1950's – Batch Processing

Pros:

- CPU kept busy, less idle time
- Monitor could provide I/O services

Cons:

- No longer interactive – longer turnaround time
- Debugging more difficult
- CPU still idle for I/O-bound jobs
- Buggy jobs could require operator intervention

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Multiprogrammed Batch Systems

- CPU is often idle
 - Even with automatic job sequencing.
 - I/O devices are slow compared to processor

Read one record from file	15 μs
Execute 100 instructions	1 μs
Write one record to file	15 μs
TOTAL	31 μs

Percent CPU Utilization = $\frac{1}{31} = 0.032 = 3.2\%$

Figure 2.4 System Utilization Example

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Uniprogramming

- Processor must wait for I/O instruction to complete before preceding

(a) Uniprogramming

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Multiprogramming

- When one job needs to wait for I/O, the processor can switch to the other job

(b) Multiprogramming with two programs

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Multiprogramming

(c) Multiprogramming with three programs


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1960's – Multiprogramming (time-sharing)

- CPU and I/O devices are multiplexed (shared) between a number of jobs
 - While one job is waiting for I/O another can use the CPU
 - SPOOLing: Simultaneous Peripheral Operation OnLine
 - 1st and simplest multiprogramming system
- Monitor (resembles O/S)
 - Starts job, spools operations, I/O, switch jobs, protection between memory

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1960's – Multiprogramming (time-sharing)



IBM System 360

Pros:

- Paging and swapping (RAM)
- Interactiveness
- Output available at completion
- CPU kept busy, less idle time

Cons:

- H/W more complex
- O/S complexity?

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1970's - Minicomputers and Microprocessors

- Trend toward many small personal computers or workstations, rather than a single mainframe.
 - Advancement of Integrated circuits
- Timesharing
 - Each user has a terminal and shares a single machine (Unix)

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1980's – Personal Computers & Networking

- Microcomputers = PC (size and \$)
- MS-DOS, GUI, Apple, Windows
- Networking: Decentralization of computing required communication
 - Not cost-effective for every user to have printer, full copy of software, etc.
 - Rise of cheap, local area networks (Ethernet), and access to wide area networks (Arpanet).

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1980's – Personal Computers & Networking

- OS issues:
 - Communication protocols, client/server paradigm
 - Data security, encryption, protection
 - Reliability, consistency, availability of distributed data
 - Heterogeneity
 - Reducing Complexity
- Ex: Byte Ordering



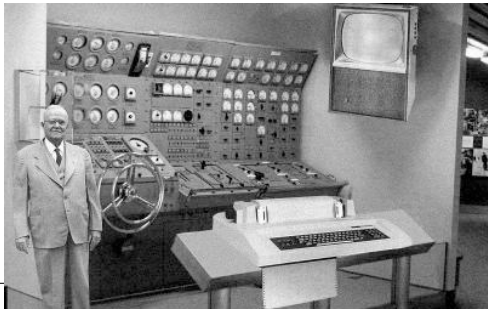
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1990's – Global Computing

- Dawn of the Internet
 - Global computing system
- Powerful CPUs cheap! Multicore systems
- High speed links
- Standard protocols (HTTP, FTP, HTML, XML, etc)
- OS Issues:
 - Communication costs dominate
 - CPU/RAM/disk speed mismatch
 - Send data to program vs. sending program to data
 - QoS guarantees
 - Security

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In the year 2000...



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2000's – Embedded and Ubiquitous Computing

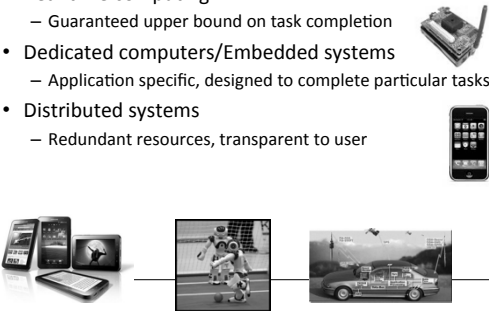
- Mobile and wearable computers
- Networked household devices
- Absorption of telephony, entertainment functions into computing systems
- OS issues:
 - Security, privacy
 - Mobility, ad-hoc networks, power management
 - Reliability, service guarantees



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2000's – Embedded and Ubiquitous Computing

- Real-time computing
 - Guaranteed upper bound on task completion
- Dedicated computers/Embedded systems
 - Application specific, designed to complete particular tasks
- Distributed systems
 - Redundant resources, transparent to user



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Multi-core

- New hotness in CPU design. Not going away.
 - Why?
- Being able to program with threads and concurrent algorithms will be a crucial job skill going forward
 - Don't leave SBU without mastering these skills
 - We will do some thread programming in Lab 3

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Editorial

- Some textbooks imply modern OSes are microkernels
- This is false
 - Windows NT and OSX were designed as microkernels
 - Then reverted to essentially monolithic designs in practice
- Linux was never a microkernel
 - Google the famous Torvalds Tanenbaum debate
- Similarly, Distributed OSes are mostly abandoned

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Object orientation

- Objects are a key feature of the Windows NT kernel design
 - IMO a good idea
- Linux actually has its own bizarre version of object orientation using C structs and function pointers
 - In Unix, everything is a file
 - How did they pull this off?
 - Poor-man's object inheritance

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Summary

- OS's began with big expensive computers used interactively by one user at a time.
- Batch systems sequences jobs to keep computer busier. Interactivity sacrificed.
- Multiprogramming developed to make more efficient use of expensive hardware and restore interactivity.
- Cheap CPU/memory/storage make communication the dominant cost.
- Multiprogramming still central for handling concurrent interaction with environment.

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<h2>Summary (2)</h2>	
<ul style="list-style-type: none">• Understand what an OS is<ul style="list-style-type: none">– Three masters– Nomenclature• Questions?	
<small>2-55</small>	