Page Frame Reclaiming

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CSE 506: Operating Systems

Logical Diagram

Today’s Lecture (kernel level mem. management)

Binary Formats
Memory Management
Interrupts
RCU
File System
Memory Allocators
Device Drivers
Scheduler
Consistency
Threads
System Calls
User
Kernel
Hardware
Memory
Disk
Net
Last time...

• We saw how you go from a file or process to the constituent memory pages making it up
  – Where in memory is page 2 of file “foo”?  
  – Or, where is address 0x1000 in process 100?

• Today, we look at reverse mapping:
  – Given physical page X, what has a reference to it?

• Then we will look at page reclamation:
  – Which page is the best candidate to reuse?
Motivation: Swapping

• Most OSes allow virtual memory to become “overcommitted”
  – Processes may allocate more virtual memory than there is physical memory in the system

• How does this work?
  – OS transparently takes some pages away and writes them to disk
  – I.e., the OS “swaps” them to disk and reassigns the physical page
Swapping, cont.

• If we swap a page out, what do we do with the old page table entries pointing to it?
  – We clear the PTE_P bit so that we get a page fault

• What do we do when we get a page fault for a swapped page?
  – We need to allocate another physical page, reread the page from disk, and re-map the new page
Choices, choices...

- The Linux kernel decides what to swap based on scanning the page descriptor table
  - Similar to the Pages array in JOS
  - I.e., primarily by looking at physical pages

- Today’s lecture:
  1) Given a physical page descriptor, how do I find all of the mappings? Remember, pages can be shared.
  2) What strategies should we follow when selecting a page to swap?
Shared memory

• Recall: A vma represents a region of a process’s virtual address space
• A vma is private to a process
• Yet physical pages can be shared
  – The pages caching libc in memory
  – Even anonymous application data pages can be shared, after a copy-on-write fork()
• So far, we have elided this issue. No longer!
Anonymous memory

- When anonymous memory is mapped, a vma is created
  - Pages are added on demand (laziness rules!)
- When the first page is added, an anon_vma structure is also created
  - vma and page descriptor point to anon_vma
  - anon_vma stores all mapping vmas in a circular linked list
- When a mapping becomes shared (e.g., COW fork), create a new VMA, link it on the anon_vma list
Example

Physical page descriptors

Process A

Virtual memory

Page Tables

Physical memory

Process B (forked)

Physical page descriptors

Page Tables
Example (2nd Page)

Physical memory

Virtual memory

Page Tables

Physical page descriptors

Process A

vma

No update?
Anonymous VMAs tend to be COW

Process B

vma

Page Tables

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Reverse mapping

• Suppose I pick a physical page X, what is it being used for?
• Many ways you could represent this
• Remember, some systems have a lot of physical memory
  – So we want to keep fixed, per-page overheads low
  – Can dynamically allocate some extra bookkeeping
Linux strategy

• Add 2 fields to each page descriptor
  • _mapcount: Tracks the number of active mappings
    – -1 == unmapped
    – 0 == single mapping (unshared)
    – 1+ == shared

• mapping: Pointer to the owning object
  – Address space (file/device) or anon_vma (process)
  – Least Significant Bit encodes the type (1 == anon_vma)
Anonymous page lookup

• Given a physical address, page descriptor index is just simple division by page size

• Given a page descriptor:
  – Look at _mapcount to see how many mappings. If 0+:
    – Read mapping to get pointer to the anon_vma
      • Be sure to check, mask out low bit

• Iterate over vmas on the anon_vma list
  – Linear scan of page table entries for each vma
    • vma-> mm -> pgdir
Example

Physical memory

Process A

Virtual memory

Page 0x10
_mapcount: 1
mapping:
(anon vma + low bit)

foreach vma

anom
vma

Process B

Page 0x10
_mapcount: 1
mapping:
(anon vma + low bit)

foreach vma

Page 0x10000
Divide by 0x1000 (4k)

Physical memory

Page Tables

Linear scan of page tables

Page Tables
File vs. anon mappings

• Given a page mapping a file, we store a pointer in its page descriptor to the inode address space
  – page->index caches the offset into the file being mapped

• Now to find all processes mapping the file...

• So, let’s just do the same thing for files as anonymous mappings, no?
  – Could just link all VMAs mapping a file into a linked list on the inode’s address_space.

• 2 complications:
Complication 1

• Not all file mappings map the entire file
  – Many map only a region of the file

• So, if I am looking for all mappings of page 4 of a file
  a linear scan of each mapping may have to filter
  vmas that don’t include page 4
Complication 2

• Intuition: anonymous mappings won’t be shared much
  – How many children won’t exec a new executable?
• In contrast, (some) mapped files will be shared a lot
  – Example: libc

• Problem: Lots of entries on the list + many that might not overlap
• Solution: Need some sort of filter
Priority Search Tree

• Idea: binary search tree that uses overlapping ranges as node keys
  – Bigger, enclosing ranges are the parents, smaller ranges are children
  – Not balanced (in Linux, some uses balance them)

• Use case: Search for all ranges that include page N

• Most of that logarithmic lookup goodness you love from tree-structured data!
retrieved. Then the algorithm visits the children (1,2,3) and (2,0,2), but it discovers
that neither of them include the page.

We won’t be able, for lack of space, to describe in detail the data structures and the
functions that implement the Linux PSTs. We’ll only mention that a node of a PST is
represented by a `prio_tree_node` data structure, which is embedded in the
`shared.prio_tree_node` field of each memory region descriptor. The
`shared.vm_set` data structure is used—as an alternative to
`shared.prio_tree_node`—to insert the memory
region descriptor in a duplicate list of a PST node. PST nodes can be inserted and
removed by executing the `vma_prio_tree_insert()` and
`vma_prio_tree_remove()` functions; both of them receive as their parameters the address of a memory region
descriptor and the address of a PST root. Queries on the PST can be performed by exe-
cuting the `vma_prio_tree_foreach` macro, which implements a loop over all memory
region descriptors that includes at least one page in a specified range of linear
addresses.

The `try_to_unmap_file()` function

The `try_to_unmap_file()` function is invoked by
`try_to_unmap()` to perform the
reverse mapping of mapped pages. This function is quite simple to describe when the
memory mapping is linear (see the section “Memory Mapping” in Chapter 16). In
this case, it performs the following actions:

1. Gets the `page->mapping->i_mmap_lock` spin lock.
2. Applies the `vma_prio_tree_foreach()` macro to the priority search tree whose
root is stored in the `page->mapping->i_mmap` field. For each
`vm_area_struct` descriptor found by the macro, the function invokes
`try_to_unmap_one()` to try
to clear the Page Table entry of the memory region that contains the page (see
the earlier section “Reverse Mapping for Anonymous Pages”). If for some reason
this function returns a `SWAP_FAIL` value, or if the
 `_mapcount` field of the page
descriptor indicates that all Page Table entries referencing the page frame have
been found, the scanning terminates immediately.

Figure 17-2
(from Understanding the Linux Kernel)

`Figure 17-2. A simple example of priority search tree`

- **Radix** – start of interval, heap = last page
- **Range is exclusive, e.g., [0, 5)**
How to find page 1?

- If in range: search both children
- If out of range: search only right or left child

Figure 17-2. A simple example of priority search tree
PST + vmas

• Each node in the PST contains a list of vmas mapping that interval
  – Only one vma for unusual mappings

• So what about duplicates (ex: all programs using libc)?
  – A very long list on the (0, filesz, filesz) node
    • I.e., the root of the tree
Reverse lookup, review

• Given a page, how do I find all mappings?
Problem 2: Reclaiming

• Until there is a problem, kernel caches and processes can go wild allocating memory

• Sometimes there is a problem, and the kernel needs to reclaim physical pages for other uses
  – Low memory, hibernation, free memory below a “goal”

• Which ones to pick?
  – Goal: Minimal performance disruption on a wide range of systems (from phones to supercomputers)
Types of pages

- Unreclaimable – free pages (obviously), pages pinned in memory by a process, temporarily locked pages, pages used for certain purposes by the kernel
- Swappable – anonymous pages, tmpfs, sharedIPC memory
- Syncable – cached disk data
- Discardable – unused pages in cache allocators
General principles

• Free harmless pages first
• Steal pages from user programs, especially those that haven’t been used recently
• When a page is reclaimed, remove all references at once
  – Removing one reference is a waste of time
• Temporal locality: get pages that haven’t been used in a while
• Laziness: Favor pages that are “cheaper” to free
  – Ex: Waiting on write back of dirty data takes time
  – Note: Dirty pages are still reclaimed, just not preferred!
Another view

• Suppose the system is bogging down because memory is scarce

• The problem is only going to go away permanently if a process can get enough memory to finish
  – Then it will free memory permanently!

• When the OS reclaims memory, we want to avoid harming progress by taking away memory a process really needs to make progress

• If possible, avoid this with educated guesses
LRU lists

- All pages are on one of 2 LRU lists: active or inactive
- Intuition: a page access causes it to be switched to the active list
  - A page that hasn’t been accessed in a while moves to the inactive list
How to detect use?

• Tag pages with “last access” time
• Obviously, explicit kernel operations (mmap, mprotect, read, etc.) can update this

• What about when a page is mapped?
  – Remember those hardware access bits in the page table?
  – Periodically clear them; if they don’t get re-set by the hardware, you can assume the page is “cold”
    • If they do get set, it is “hot”
Big picture

• Kernel keeps a heuristic “target” of free pages
  – Makes a best effort to maintain that target; can fail

• Kernel gets really worried when allocations start failing
  – In the worst case, starts out-of-memory (OOM) killing processes until memory can be reclaimed
Editorial

• Choosing the “right” pages to free is a problem without a lot of good science behind it
  – Many systems don’t cope well with low-memory conditions
  – But they need to get better
    • (Think phones and other small devices)

• Important problem – perhaps an opportunity?
Summary

• Reverse mappings for shared:
  – Anonymous pages
  – File-mapping pages

• Basic tricks of page frame reclaiming
  – LRU lists
  – Free cheapest pages first
  – Unmap all at once
  – Etc.